# University of California - Berkeley <br> Department of Economics ECON 201A Economic Theory <br> Choice Theory 

Fall 2023

## Preferences, Utility, and Choice (Kreps Ch. 1 and Rubinstein Ch. 1-3)

Aug 31 and Sep 5

## Ever since Ellsberg...

The Senator Gravel Edition

## Pentagon Papers

The Defense Department History of United States Decisionmaking on
VIETNAM
VOLUME ONE


Wive Che Washington Jost


## Experiments à la Ellsberg

Consider the following four two-color Ellsberg-type urns (Halevy, 2007):
I. 5 red balls and 5 black balls
II. an unknown number of red and black balls
III. a bag containing 11 tickets with the numbers $0-10$; the number written on the drawn ticket determines the number of red balls
IV. a bag containing 2 tickets with the numbers 0 and 10 ; the number written on the drawn ticket determines the number of red balls

## Kreps' word of advice...

A recommendation for this course which extends to any microeconomics book/paper (mathematical in character):
$\Longrightarrow$ Read carefully and slowly for details - follow the details of the proofs one step at a time, have a pad and pen, make notes, and finish arguments.
$\Longrightarrow$ But do not lose the "plot line," which is just as important - What is the framework? What are the results? How do the results tie together?

Constructing (mathematical) proofs is a skill you learn best - and perhaps only - by doing.

## Outline

Preferences

The preferences of the $\mathcal{D} \mathcal{M}$ for $x, y \in X$ (a finite set of objects of choice) are specified by a binary relation $\succsim$ where $x \succsim y$ is read " $x$ is at least as good as $y$ " or as " $x$ is weakly preferred to $y$."

- $\succsim$ on $X$ is complete if for $\forall x, y \in X$, either $x \succsim y$ or $y \succsim x$ (or both).
- $\succsim$ on $X$ is is transitive if $x \succsim y$ and $y \succsim z$ then $x \succsim z$.
$\Longrightarrow$ It would be nice to begin with something less abstract / more concrete but $\succsim$ are the logical starting point for choice theory...


## $\underline{\text { Utility }}$

The $\mathcal{D M}$ 's utility function is a real-valued function $u: X \rightarrow \mathbb{R}$. It represents $\succsim$ if

$$
\begin{equation*}
x \succsim y \Longleftrightarrow u(x) \geq u(y) \tag{1}
\end{equation*}
$$

that is, the $\mathcal{D} \mathcal{M}$ regards items of higher utility as better.
$\Longrightarrow$ Many different utility functions represents the same preferences so utility numbers (for now) have only ordinal and not cardinal meaning.

## Choice

A choice function $c$ such that $c(A) \subseteq A$ which specifies for each nonempty subset $A \subseteq X$ what the $\mathcal{D} \mathcal{M}$ would be 'content' to have.

- $c$ is generated by $\succsim$ if $\forall A$

$$
\begin{equation*}
c(A)=\{x \in A: x \succsim y \forall y \in A\} . \tag{2}
\end{equation*}
$$

- $c$ is generated by utility maximization of $u$ if $\forall A$

$$
\begin{equation*}
c(A)=\{x \in A: u(x) \geq u(y) \forall y \in A\} \tag{3}
\end{equation*}
$$

Properties of choice functions
C1 A choice function $c$ satisfies finite nonemptiness if $c(A) \neq \emptyset$ for $\forall A \in$ $\mathcal{A}$ where $\mathcal{A}$ is all the nonempty subsets $A \subseteq X$.

C2 A choice function $c$ satisfies choice coherence if $\forall x, y \in X$ and $\forall A, B \in \mathcal{A}$, if $x, y \in A \cap B, x \in c(A)$, and $y \notin c(A)$, then $y \notin c(B)$.

The main result(s) ( $X$ finite)

- If $c$ satisfies $C 1$ and C2, then there exist both a complete and transitive $\succsim$ and $u$ that 'produce' choices according to $c$ via (2) and (3), respectively.
- If $\succsim$ is complete and transitive, then the $c$ it produces via (2) satisfies C1 and C2, and there exists $u$ that represents $\succsim$ via (1).
- Given any $u$, the $c$ it produces via (3) satisfies $C 1$ and $C 2$ and the $\succsim$ it represents via (1) is complete and transitive. And the $c$ produced by that $\succsim$ via (2) is precisely the $c$ produced from $u$ via (3).
- In words, choice behavior that satisfies C1 and C2 is equivalent to complete and transitive $\succsim$-optimality and both of which are equivalent to $u$ maximization.
- This conglomerate (the two pairs of assumptions) is the standard model of choice in microeconomics:
$\rightarrow \quad$ prove and generalize
$\rightarrow$ extend to infinite $X$
$\rightarrow$ comments, extensions, variations, and criticisms...
$\Longrightarrow$ Well, we know where we're goin' But we don't know where we've been... (Talking Heads |Little Creatures)


## Preferences

We start by taking up the following (simple) story:

- There is a (finite) set $X$ of options and a $\mathcal{D} \mathcal{M}$ who is willing to express her/his preferences among these options by making paired comparisons.
- We want to 'fully' describe the preferences of the $\mathcal{D} \mathcal{M}$ toward the options in $X$, thinking about preferences independently of choice.
$\Longrightarrow$ Your attitude toward the PhD programs you have applied to before finding out to which of them you have been accepted...


## Rubinstein's (imaginary) questionnaires

Questionnaire $Q(x, y)$ : Check one of the following three boxes:
$\square \quad$ I prefer $x$ to $y(x \succ y)$
$\square \quad$ I prefer $y$ to $x(y \succ x)$
$\square \quad$ I am indifferent ( $I$ )

- I cannot compare, do not know, have no opinion, prefer both, need to consult a friend (and so on) are "illegal" answers...
- We use the symbol $I$ for indifference because

$$
x \sim y \Longleftrightarrow x \succsim y \text { and } y \succsim x
$$

For "legal" answers to the questionnaire $Q$ to qualify as preferences must satisfy two "consistency" requirements:

1. No "framing effect" - the answer to $Q(x, y)$ must be the answer to $Q(y, x)$.
2. The answers to $Q(x, y)$ and $Q(y, z)$ are consistent (in the transitive sense...) with the answer to $Q(x, z)$.

And the answer to the "silly" question $Q(x, x)$ must be I am indifferent $(I)$ but we will consider only distinct options $x \neq y$.

Preferences (based on questionnaire $Q$ ): Preferences are a function $f$ that assigns to any $x, y \in X$ one of the three "values" $x \succ y, y \succ x$, or $I$ so that for any $x, y, z \in X$, the following two properties hold:

1. No order effect: $f(x, y)=f(y, x)$.
2. Transitivity:

$$
\begin{aligned}
& f(x, y)=x \succ y \text { and } f(y, z)=y \succ z \Longrightarrow f(x, z)=x \succ z \\
& f(x, y)=I \text { and } f(y, z)=I \Longrightarrow f(x, z)=I
\end{aligned}
$$

All other consistency requirements such as $f(x, y)=x \succ y$ and $f(y, z)=$ $I \Longrightarrow f(x, z)=x \succ z$ follow from the above conditions.

Questionnaire $R(x, y)$ : Is $x \succsim y$ ? Check one (and only one) of the following two boxes:
$\square \quad$ Yes
$\square \quad$ No

If the answer to $R(x, y)$ is "Yes" then we identify the response to this with $x \succsim y$. To qualify as preferences, responses must satisfy two conditions:

1. At least one answer to $R(x, y)$ and $R(y, x)$ as well as the answer to $R(x, x)$ must be "Yes."
2. For any $x, y, z \in X$, if the answers to $R(x, y)$ and $R(y, z)$ are "Yes" then so is the answer to $R(x, z)$.

Preferences (based on questionnaire $R$ ): Preferences are a binary relation $\succsim$ satisfying:

1. Reflexivity: $\forall x \in X, x \succsim x$.
2. Completeness: $\forall x, y \in X, x \succsim y$ and/or $y \succsim x$.
3. Transitivity: $\forall x, y, z \in X$, if $x \succsim y$ and $y \succsim z$ then $x \succsim z$.

This is the conventional definition of preferences but the two definitions of preferences are equivalent.

The equivalence of the two definitions of preferences

| $Q(x, y)$ and $Q(y, x)$ | $R(x, y)$ and $R(y, x)$ |
| :---: | :---: |
| $x \succ y$ | Yes No |
| $y \succ x$ | No Yes |
| $I$ | Yes Yes |

- Equivalence of two definitions (in microeconomics) requires the existence of a one-to-one (bijection) correspondence that preserves the interpretation...
- Exercise: Construct a "translation" between answers to $Q$ (qualify as preferences by the 1 st definition) and answers to $R$ (qualify as preferences by the 2nd definition).


## Sources of intransitivity

(1) Aggregating of considerations:

- Let $X=\{a, b, c\}$ and assume the $\mathcal{D M}$ 's attitude is $x \succ y$ if the majority of considerations supports $x$.
- The aggregation of three considerations such that

$$
a \succ_{1} b \succ_{1} c \quad b \succ_{2} c \succ_{2} a \quad \text { and } \quad c \succ_{3} a \succ_{3} b
$$

leads to violating transitivity.
(2) Using similarities:

- Let let $X=\mathbb{R}$ and assume the $\mathcal{D} \mathcal{M}$ 's attitude is "the larger the better" but $\mathcal{D} \mathcal{M}$ is unable to determine whether $a>b$ unless $|a-b| \geq 1$, thus

$$
x \succ y \text { if } x \geq y+1
$$

- This is not a preference relation because

$$
0 \sim \frac{2}{3} \text { and } \frac{2}{3} \sim \frac{4}{3} \quad \text { but not } 0 \sim \frac{4}{3}
$$

## Utility

It is possible to avoid the notion of a utility representation and to "do economics" solely based on the notion of preferences, but...
$\Longrightarrow u: X \rightarrow \mathbb{R}$ represents the preference $x \succsim y$ if $\forall x, y \in X$

$$
x \succsim y \Longleftrightarrow u(x) \geq u(y)
$$

If the function $u$ represents the preference relation $\succsim$, we refer to it as a utility function and say that $\succsim$ has a utility representation.

If $u$ represents $\succsim$, then for any strictly increasing function $f: \mathbb{R} \rightarrow \mathbb{R}$, the function $v(x)=f(u(x))$ also represents $\succsim$.

$$
\begin{aligned}
& a \succsim b \\
& \text { ॥ } \\
& u(a) \geq u(b) \\
& \text { ॥ } \\
& f(u(a)) \underset{\substack{\mathbb{1}}}{\geq} f(u(b)) \\
& v(a) \geq v(b)
\end{aligned}
$$

If $\succsim$ has a utility representation, then it has an infinite number of such representations...

## Existence of a utility representation

$X$ is finite:

- If $\succsim$ is a preference relation on a finite set $X$, then $\succsim$ has a utility representation with values that are natural numbers.
$\underline{X}$ is countable (if infinite has a one-to-one correspondence with the set of natural numbers):
- If $\succsim$ is a preference relation on a countable set $X$, then $\succsim$ has a utility representation with a range $(0,1)$.

Proof (for finite $X$ ):

First we show by induction that any finite set $A \subseteq X$, has a mini$\mathrm{mal} /$ maximal element $-a \in X$ is minimal if $a \precsim x$ for any $x \in X$.

- If $A$ is a singleton, then by reflexivity its (single) element is minimal.
- By the inductive assumption, the set $A-\{x\}$ for some $x \in A$ has a minimal element, say $y$.
- If $x \succsim y$, then $y$ is minimal in $A$. If $y \succsim x$, then by transitivity $z \succsim x$ for all $z \in A-\{x\}$ so $x$ is minimal.

Next let $X_{k+1}$ be set of minimal elements in $X-X_{1}-\cdots-X_{k}$ and note that $X_{k+1} \neq \emptyset$.

- Let $u(x)=k$ if $x \in X_{k}$ ( $x$ is "eliminated" after $k$ steps) and note that because $X$ is finite $u(x) \leq|X|$.
- If $a \succ b$ then $a \notin X_{1} \cup X_{2} \cup \cdots \cup X_{u(b)}$ so $u(a)>u(b)$ and if $a \sim b$ then $u(a)=u(b)$.
$\Longrightarrow$ Hence, $u$ represents $\succsim$. The proof of the case where $X$ is countable is left as an exercise. No further assumptions on $\succsim$ are needed.


## Lexicographic preferences

"...even though they are derived from a simple and commonly used procedure," lexicographic preferences do not have a utility representation by $u: X \rightarrow$ $\mathbb{R}$.

Let $X=[0,1] \times[0,1]$ (the unit square) and let $x \succsim_{k} y$ if $x_{k} \geq y_{k}$. The lexicographic preferences $\succsim_{L}$ induced by $\succsim_{1}$ and $\succsim_{2}$ are given by

$$
\left(a_{1}, a_{2}\right) \succsim_{L}\left(b_{1}, b_{2}\right) \text { if } a_{1}>b_{1} \text { or } a_{1}=b_{1} \text { and } a_{2} \geq b_{2}
$$

Proof: Assume $u: X \rightarrow \mathbb{R}$ represents $\succsim_{L}$ toward contradiction.

- For any $a \in[0,1],(a, 1) \succ(a, 0)$ and thus $u(a, 1)>u(a, 0)$. Let $q(a)$ be a rational number in the nonempty interval

$$
I_{a}=(u(a, 0), u(a, 1))
$$

$-q(a):[0,1] \rightarrow \mathbb{R}$ and it is one-to-one function since

$$
b>a \Longrightarrow(b, 0) \succ_{L}(a, 1) \Longrightarrow u(b, 0)>u(a, 1)
$$

- It follows that the intervals $I_{a}$ and $I_{b}$ are are disjoint so $q(a) \neq q(b)$ but the cardinality of the rational numbers is lower than that of the continuum.


## Continuity of preferences

Continuity condition guarantees the existence of a utility representation when $X$ is an infinite subset of a Euclidean space (Debreu's theorem).
I. $\succsim$ on $X$ is continuous if whenever $a \succ b$, there are (small) balls $B_{a}$ and $B_{b}$ around $a$ and $b$, respectively, such that $a \succ b$ for all $x \in B_{a}$ and $y \in B_{b}$.
II. $\succsim$ on $X$ is continuous if $\left\{\left(a_{n}, b_{n}\right)\right\}$ a sequence of pairs satisfying $a_{n} \succsim b_{n}$ for all $n$ and $a_{n} \rightarrow a$ and $b_{n} \rightarrow b$, then $a \succsim b$. That is, if the graph of $\succsim$

$$
\{(x, y) \mid x \succsim y \subseteq X \times X\}
$$

is closed.

Three remarks:

1. $\succsim$ on $X$ satisfies (I) iff it satisfies (II).
2. If $\succsim$ can be represented by a continuous $u$ then it is continuous.
3. $\succsim_{L}$ is not continuous: for "small" $\varepsilon$

$$
(1,1) \succ_{L}(1,0) \text { but }(1-\varepsilon, 1) \prec_{L}(1,0) .
$$

$\Longrightarrow$ Prove (1) and (2) as an exercise.


## Debreu's $(1954,1960)$ theorem

One of the classical results in economic theory:
Debreu's theorem Let $\succsim$ be a continuous preference relation on $X$, which is a convex subset of $\mathbb{R}^{n}$. Then $\succsim$ has a continuous utility representation.

The proof (in Rubinstein), relies on the mathematical concept of a dense set - a set $Y$ is said to be dense in $X$ if every non-empty open set $B \subset X$ contains an element in $Y$. Any set $X \subseteq \mathbb{R}^{n}$ has a countable dense subset.

## Choice

A choice function $c$ such that $c(A) \subseteq A$ which specifies for each nonempty subset $A \subseteq X$ what the $\mathcal{D M}$ would be 'content' to have.

- Not all choice problems are always relevant. We thus allow the $\mathcal{D M}$ 's behavior to be defined only on a restricted set $D$ of subsets of $X$ and refer to a pair $(X, D)$ as context.
- An example: the $\mathcal{D M}$ is choosing whether to remain with the status quo $s$ or choose an element in a set $Y$. Than $X=Y \cup\{s\}$ and $D$ is the set of all subsets of $X$ that contain $s$.

A rational choice function:

- A rational $\mathcal{D} \mathcal{M}$ has a preference relation $\succsim$ on the set $X$ and and given any choice problem $A \subseteq X, c(A)$ is $\succsim$-optimal.
- The induced choice function $c_{\succsim}$ is the function that assigns to every nonempty set $A \in D$ the $\succsim$-best element in $A$.

A choice function $c$ can be rationalized if there is a preference relation $\succsim$ on $X$ such that $c=c_{\succsim}$, that is

$$
c(A)=c_{\succsim}(A) \text { for any } A \subseteq X
$$

Rubinstein's condition $\alpha$ : a choice function $c$ satisfies condition $\alpha$ if for any two problems $A, B \in D$, if $A \subset B$ and $c(B) \in A$, then $c(A)=c(B)$.

- The 2 nd-best choice function violates condition $\alpha-c(A)$ is the alternative from $A$ that is the $\succsim$-maximal from among the non-maximal alternatives...
- It is a sufficient condition for $c$ to be formulated as-if the $\mathcal{D M}$ is maximizing some preference relation $\succsim$.
I. Let $c$ be a choice function defined on a domain containing at least all subsets of $X$ of size of at most 3 . If $c$ satisfies condition $\alpha$, then there is a preference $\succsim$ on $X$ such that $c=c_{\succsim}$.
II. Let $c$ be a choice function defined on a domain $D$ satisfying that if two problems $A, B \in D$ then $A \cup B \in D$. If $c$ satisfies condition $\alpha$, then there is a preference $\succsim$ on $X$ such that $c=c_{\succsim}$.

We will prove (II) and prove (I) when we consider choice correspondences (next week).

Proof: Define a binary relation as $x R y$ if $x \neq y$ and there is a choice problem $A \in D$ such that $y \in A$ and $c(A)=x$.

- If $x R y$ and $y R z$ then there is $A \in D$ such that $y \in A$ and $c(A)=x$, and $B \in D$ such that $z \in B$ and $c(B)=y$.
- The set $A \cup B \in D$ and $c(A \cup B)$ is in either $A$ or $B$ and thus it is either $x$ or $y$ (by condition $\alpha$ ).
- But if $c(A \cup B)=y$ then $c(A \cup B)=c(A)=y$ so $c(A \cup B)=x$ (again by condition $\alpha$ ) and thus $x R z$.
! An acyclic and asymmetric relation $R$ extends to a preference relation $\succsim$ (to be proved later) so $c(A) \succsim x$ for all $x \in A$ and thus $c(A)=c_{\succsim}(A)$.


## Choice correspondences

The choice function $c_{\succsim}(A)$ requires that a single maximal element each choice problem - undefined for a preference relation with indifferences.

A choice correspondence $c$ assigns to every nonempty $A \in D$ a nonempty subset of $A$ so

$$
\emptyset \neq c(A) \subseteq A
$$

$\Longrightarrow$ One (behavioral) interpretation of $c(A)$ is the set of all elements in $A$ that are satisfactory in the sense of an "internal equilibrium."
$\Longrightarrow$ Another (behavioral) interpretation is that $c(A)$ is the set of all elements in $A$ that are can be satisfactory under some possible circumstances not included in the description of the set $A$.

- Let $(A, f)$ be an extended choice set where $f$ is the 'frame' that accompanies the set $A$.
- Then $c(A, f)$ is the choice of the $\mathcal{D} \mathcal{M}$ from $A$ given frame $f$

$$
C(A)=\{x \mid x=c(A, f) \text { for some } f\}
$$

Two properties of choice correspondences:
( $\alpha$ ) If $\alpha \in A \subset B$ and $\alpha \in c(B) \Longrightarrow \alpha \in c(A)$.
( $\beta$ ) If $\alpha, b \in A \subset B, \alpha, b \in c(A)$ and $b \in c(B) \Longrightarrow \alpha \in c(B)$.

For any domain $D$ such that if $A, B \in D$ then $A \cap B \in D, \alpha$ and $\beta$ are equivalent to the weak axiom of revealed preference (below).
$\Longrightarrow$ Verify this, but before provide examples of choice correspondences that satisfy one property but not the other...

The weak axiom of revealed preference (WA):
$c$ satisfies the WA if whenever $x, y \in A \cap B$

$$
x \in c(A) \text { and } y \in c(B) \Longrightarrow x \in c(B)
$$

In words, if $x$ is revealed to be at least as good as $y$, then $y$ is not revealed to be strictly better than $x$. Condition $\alpha$ for choice functions is replaced for correspondences by the WA.

Let $c$ be a choice correspondence defined on a domain containing at least all subsets of $X$ of size of at most 3. If $c$ satisfies properties $\alpha$ and $\beta$, then there is a preference $\succsim$ on $X$ such that $c=c_{\succsim}$.

Proof: Define $x \succsim y$ if $x \in c(\{x, y\})$. The relation $\succsim$ is a preference:

$$
\begin{array}{ll}
\text { Reflexive } & c(\{x\})=x \\
\text { Complete } & c(\{x, y\}) \neq \emptyset
\end{array}
$$

To see that is also transitive note that if $x \succsim y$ and $y \succsim z$ then $x \in$ $c(\{x, y\})$ and $y \in c(\{y, z\})$.

- If $y \in c(\{x, y, z\})$ resp. $z \in c(\{x, y, z\})$ then

$$
\begin{aligned}
\alpha & \Longrightarrow y \in c(\{x, y\}) \\
\beta & \Longrightarrow x \in c(\{x, y, z\})
\end{aligned}
$$

resp.

$$
\begin{aligned}
\alpha & \Longrightarrow z \in c(\{y, z\}) \\
\beta & \Longrightarrow y \in c(\{x, y, z\})
\end{aligned}
$$

Thus, in all cases $x \in c(\{x, y, z\})$ so $\alpha \Longrightarrow x \in c(\{x, z\})$ and therefore $x \succsim z$.

- It remains to be shown that $c(B)=c_{\succsim}(B)$ :
$\rightarrow x \in c(B)$ then (by $\alpha) x \in c(\{x, y\})$ for every $y \in B$ so

$$
x \succsim y \Longrightarrow x \in c_{\succsim}(B)
$$

$\leftarrow x \in c_{\succsim}(B)$ and $y \in c(B)$ then $x \in c(\{x, y\})$ since $x \succsim y$

$$
\begin{aligned}
\alpha & \Longrightarrow y \in c(\{x, y\}) \\
\beta & \Longrightarrow x \in c(B) .
\end{aligned}
$$

## Last word: procedural rationality

Consider the following Herbert Simon's so-called satisficing procedure:
$-v: X \rightarrow \mathbb{R}$ - the valuation of the elements in $X$ where $v^{*}$ is a threshold of satisfaction.

- Given a set $A \subseteq X, L(A, O)$ is the $\mathcal{D} \mathcal{M}$ 's list according to her/his ordering $O$.
$c(A)$ is the first element in $L(A, O)$ that has a $v$-value at least as large as $v^{*}$ (and if there is no such element, the last element in $L(A, O)$ is chosen).
$\Longrightarrow$ The choice function $c(A)$ induced by this procedure satisfies condition $\alpha \ldots$

